

Recap

HW 3 posted

Due: 2/21/25

I-Projections: Closed, convex  $\Pi$

$$P^* = \operatorname{argmin}_{P \in \Pi} D(P \parallel Q) = \operatorname{Proj}_{\Pi}(Q)$$

Defect inequality:  $D(P \parallel Q) \geq D(P \parallel P^*) + D(P^* \parallel Q)$

Linear families:  $\mathcal{L} = \{P \mid \sum_{x \sim P} f_i(x) = \alpha_i \quad \forall i \in [K]\}$

$$D(P \parallel Q) = D(P \parallel P^*) + D(P^* \parallel Q)$$

$$P^*(x) = C \cdot Q(x) \cdot e^{\sum \lambda_i f_i(x)}$$

# Matrix Scaling (Sinkhorn)

[Wigderson]

Operator Scaling

$$\triangleright M \in \mathbb{R}_+^{n \times n} \quad M_{ij} > 0 \quad \forall i, j$$

There exist diagonal matrices  $D_1, D_2$  st.  $M' = D_1 M D_2$  is doubly stochastic

i.e.

$$\forall i \quad \sum_j M'_{ij} = 1 \quad \forall j \quad \sum_i M'_{ij} = 1$$

Proof:

$$M \succeq Q$$

doubly stochastic  $\Rightarrow d$

$$B = \sum_j M_{ij}$$

$$M'' = \frac{M}{B}$$

Form of the optimizer

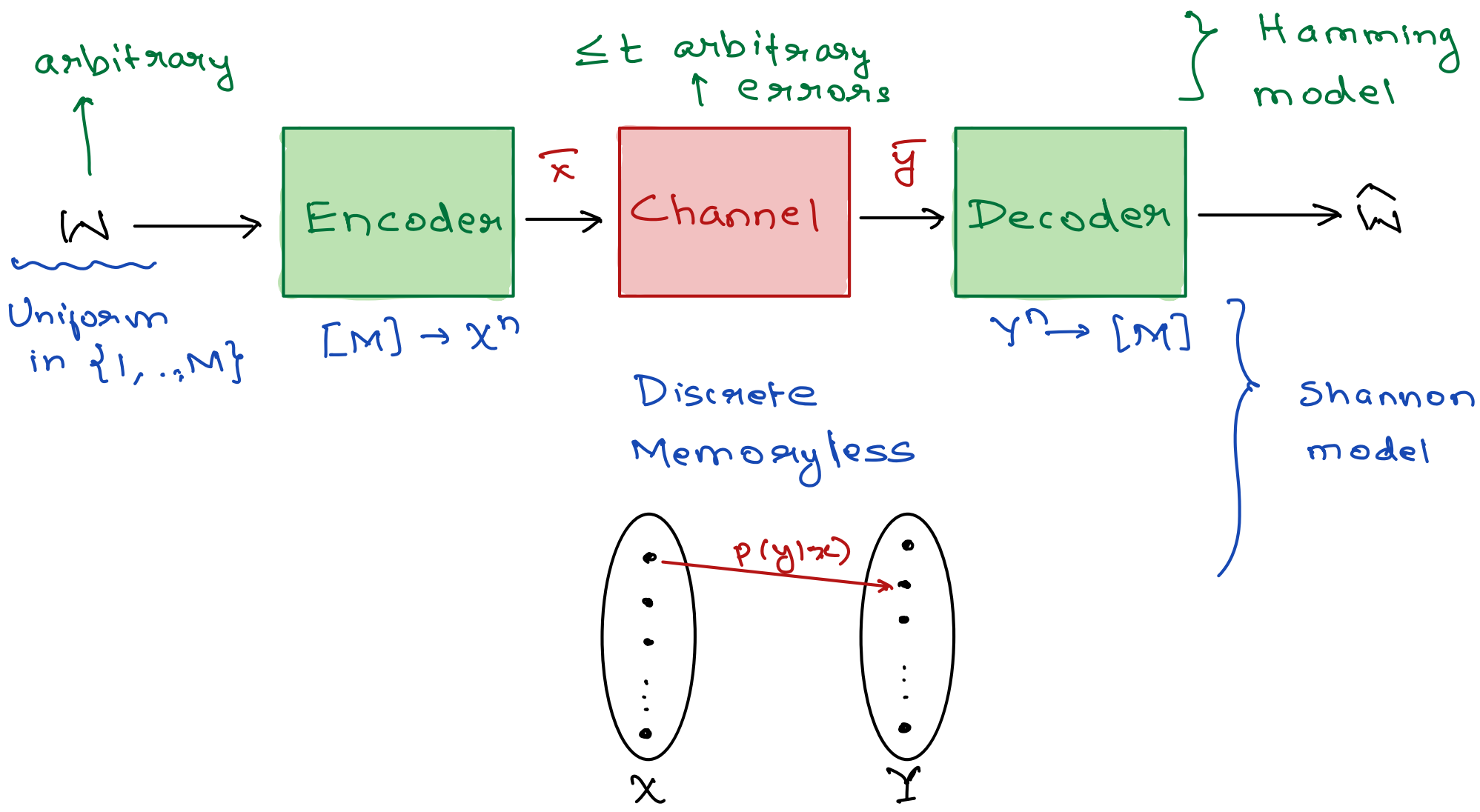
$$P^*(i,j) = c_0 \cdot Q(i,j) \cdot \exp\left(\sum_{i_0} \lambda_{i_0} f_{i_0}(i,j) + \sum_{j_0} \mu_{j_0} g_{j_0}(i,j)\right)$$

$$P^*(i,j) = c_0 \cdot M_{ij}^1 \cdot e^{\lambda_i + \mu_j}$$

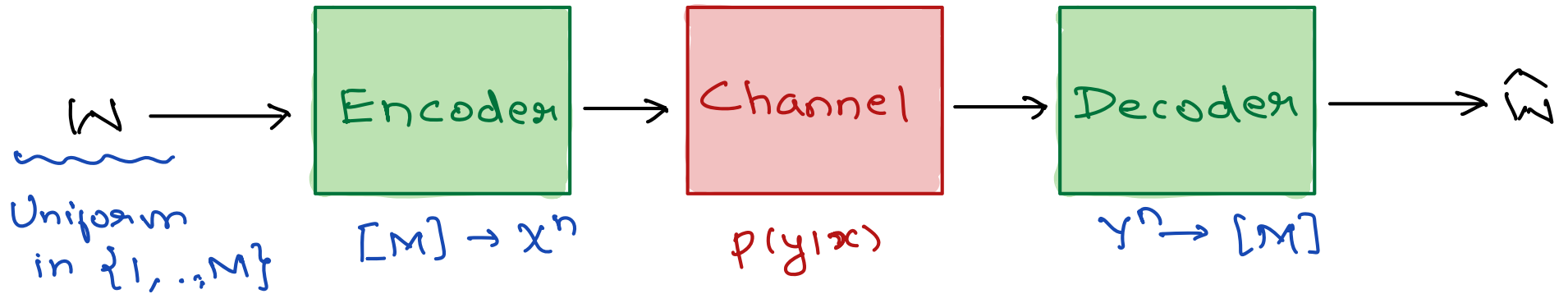
$$= \underbrace{\sqrt{c_0} \cdot e^{\lambda_i}}_{d_i} \cdot M_{ij}'' \cdot \underbrace{\sqrt{c_0} \cdot e^{\mu_j}}_{d_j'}$$

$$= (D_1 M'' D_2)_{ij}$$

# Error Correcting Codes



# The Shannon model



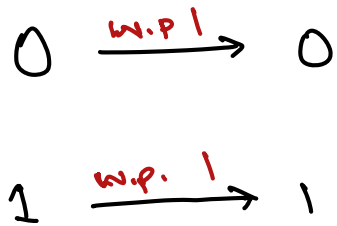
Rate: 
$$R = \frac{\log M}{\log |\mathcal{X}|^n} = \frac{\log M}{n \cdot \log |\mathcal{X}|}$$

Capacity: 
$$C = \max_{P(x)} I(x; Y)$$

# Channels and Capacities ( $X = \{0, 1\}$ )

$$C = \max_{P(X)} I(X; Y)$$

## Noiseless

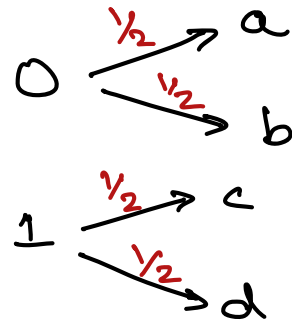


$$\max_{P(X)} I(X; Y)$$

$$= \max_{P(X)} H(X)$$

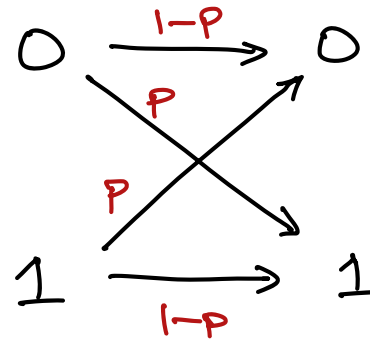
$$= 1$$

## Noiseless, random



$$= 1$$

## Binary Symmetric

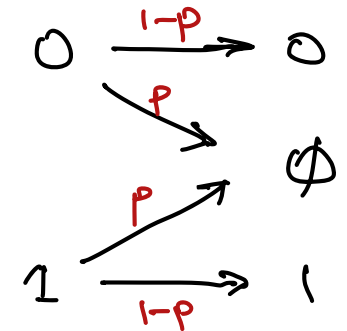


$$I(X; Y)$$

$$= H(X) - H_2(p)$$

$$C = 1 - H_2(p)$$

## Binary Erasure

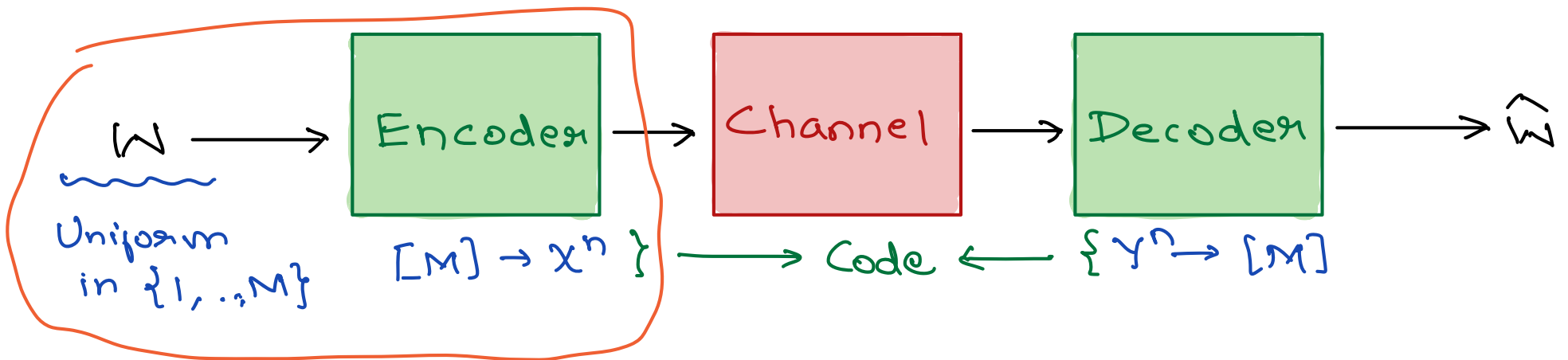


$$H(X; Y) = p \cdot H(X)$$

$$I(X; Y) = (1-p) H(X)$$

$$C = 1-p$$

# Channel Coding Theorem



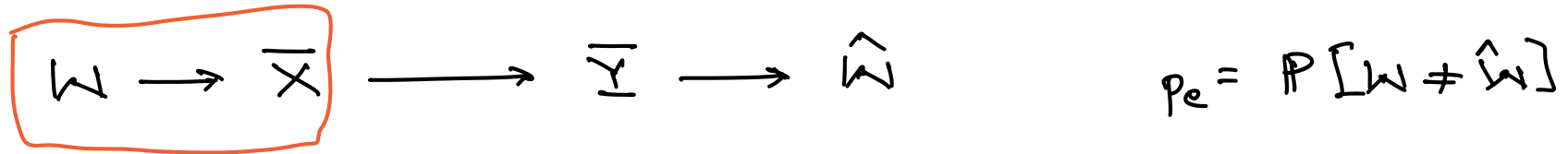
-  $P_e = P[W \neq \hat{W}]$

$\log(M) = n \cdot R$   
 $M = 2^{n \cdot R}$

- Achievable rate  $R$ :  $\exists$  Sequence of codes, rate  $\geq R$   
 $P_e^{(n)} \rightarrow 0$  as  $n \rightarrow \infty$

►  $\sup \{R \mid R \text{ is achievable}\} = R^* = C = \max_{P(x)} I(X; Y)$

Achievable  $R \leq C$



Fano:  $H_2(p_e) + p_e \cdot \underbrace{\log M}_{nR} \geq H(W | \hat{W})$

$\Rightarrow \frac{H(W)}{nR} - I(W; \hat{W})$

$\Rightarrow R - I(\bar{X}; \hat{W})$

$$I(\bar{x}; \bar{y}) = H(\bar{y}) - H(\bar{y} | \bar{x})$$

$$= \sum_i \underbrace{H(y_i | y_1 \dots y_{i-1})}_{= H(y_i)} - \sum_i \underbrace{H(y_i | y_1 \dots y_{i-1} \bar{x})}_{= H(y_i | x_i)}$$

$$\leq \sum_i I(x_i; y_i)$$

$$\leq n \cdot C$$

$$1 + p_e \cdot nR \geq nR - n \cdot C$$

$$\underbrace{\frac{1}{n}}_{\rightarrow 0} + \underbrace{p_e \cdot R}_{\rightarrow 0} + \dots \geq R$$