

Markov Chains

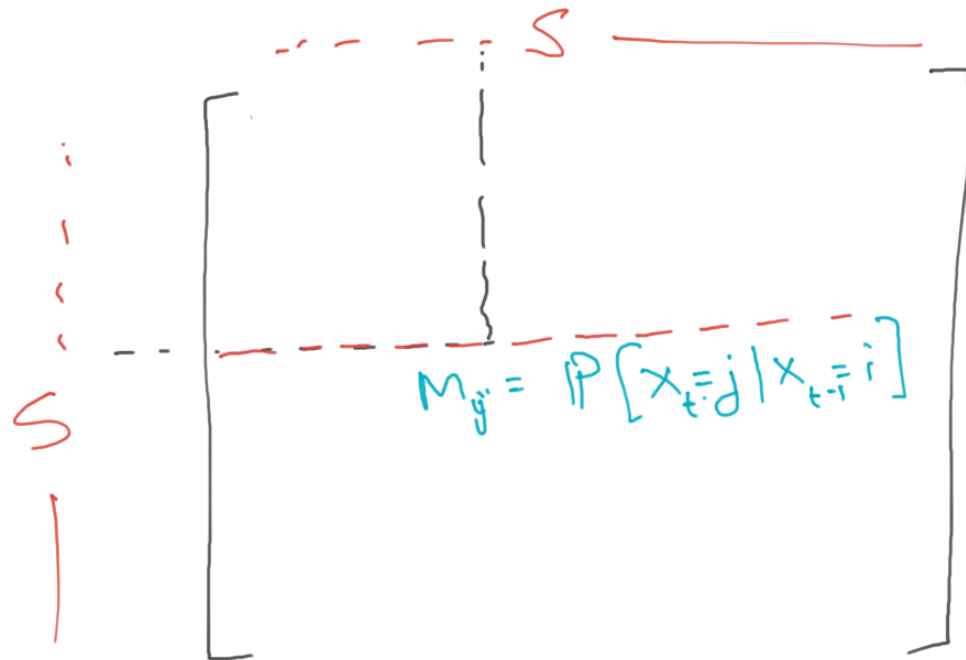
$X_0 \quad X_1 \quad X_2 \quad \dots \dots \dots$

$$P[X_t = a_t \mid X_{t-1} = a_{t-1}, \dots, X_0 = a_0] = P[X_t = a_t \mid X_{t-1} = a_{t-1}]$$

Random processes with limited history

$X_0, X_1, \dots \in S$ (finite) state space
 $|S| = n$

$$P[X_t = j \mid X_{t-1} = i] = P[X_{t'} = j \mid X_{t'-1} = i] \left. \vphantom{P[X_t = j \mid X_{t-1} = i]} \right\} \begin{array}{l} \text{time} \\ \text{homogeneous} \end{array}$$



Transition Matrix M

$$- M_{ij} \geq 0 \quad \forall i, j$$

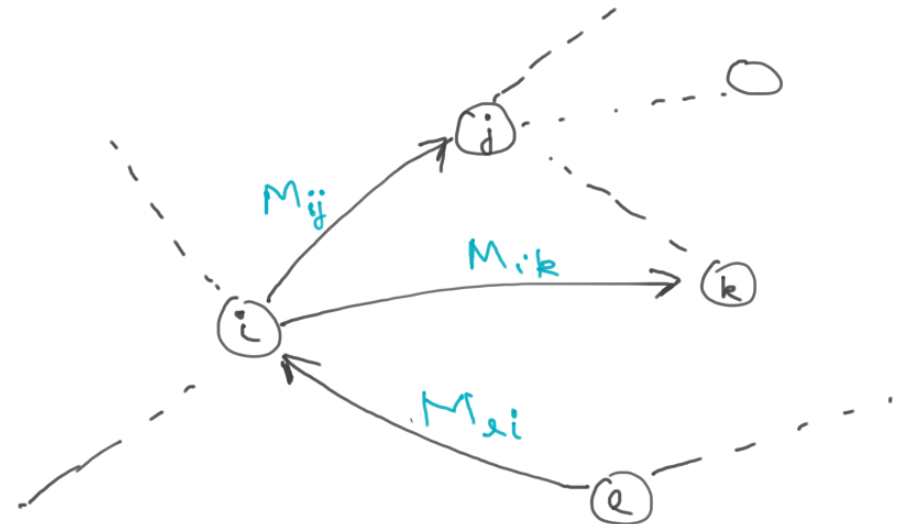
$$- \sum_j M_{ij} = 1 \quad \forall i$$

stochastic matrix

Markov Chains as directed graphs

$$M_{ij} = P[x_t = j | x_{t-1} = i]$$

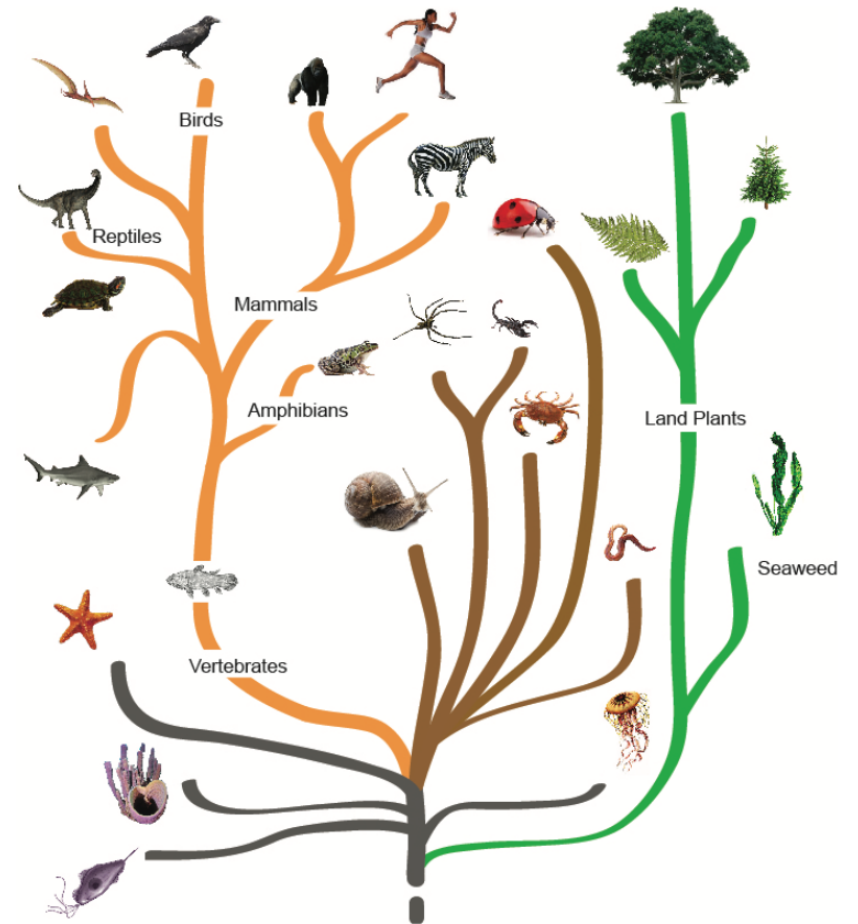
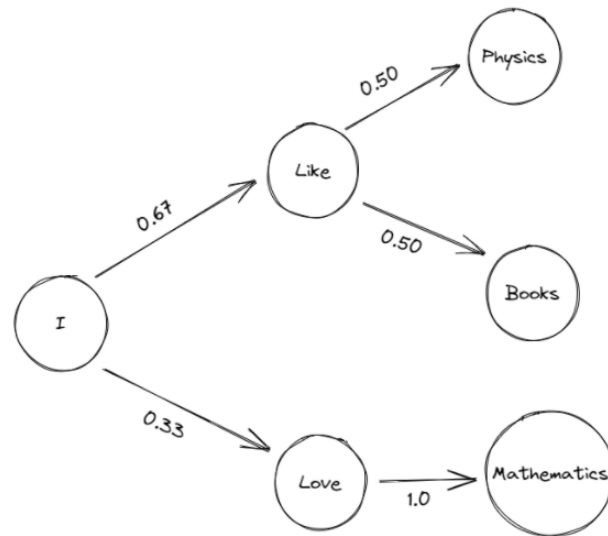
$\equiv$



Vertex Set = S

Ex: $P[x_t = j | x_0 = i] = (M^t)_{ij}$

Examples (maybe)



Examples



Card Shuffling

$$x_{t+1} = \sqrt{1-\beta} \cdot x_t + \sqrt{\beta} g$$

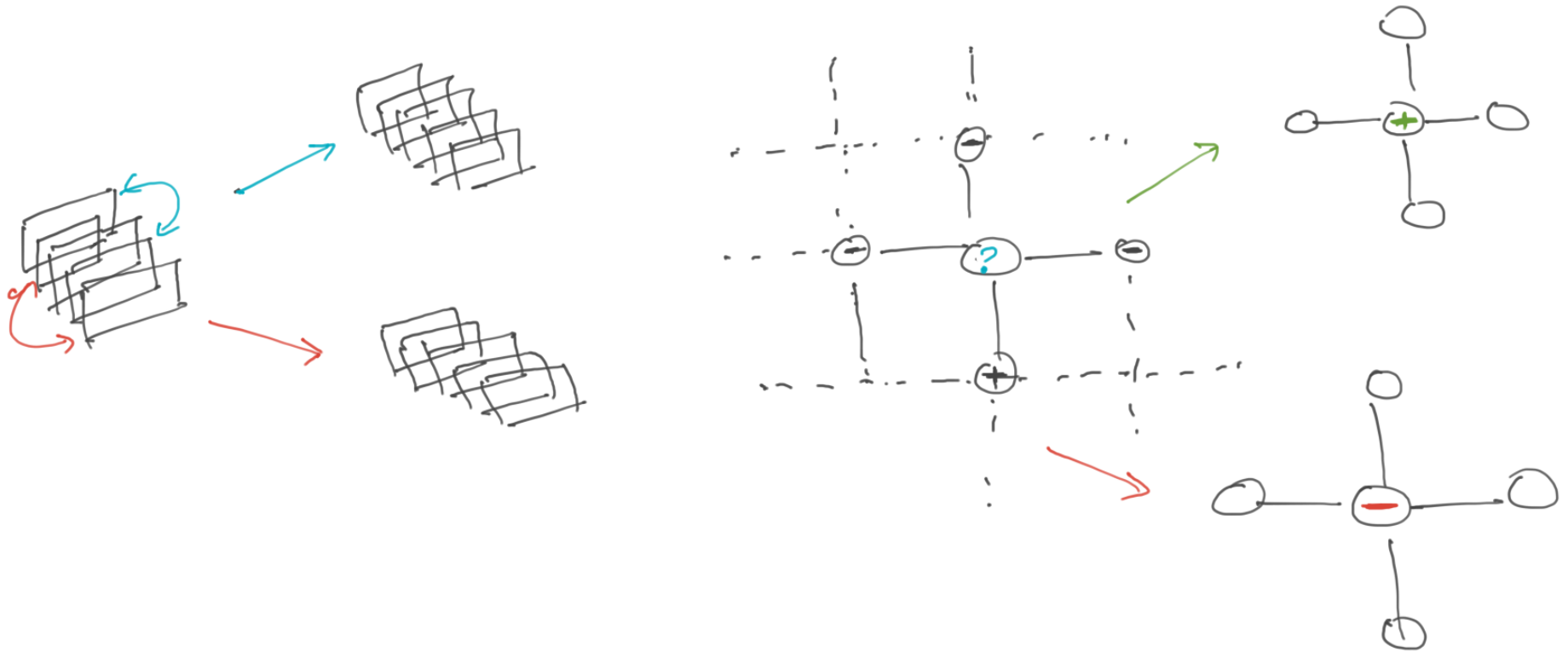
$$g \sim N(0, 1)$$

Gaussian noise

$$P[x_1, x_2, \dots] = \frac{e^{-\beta H(x)}}{Z}$$

Ising Model

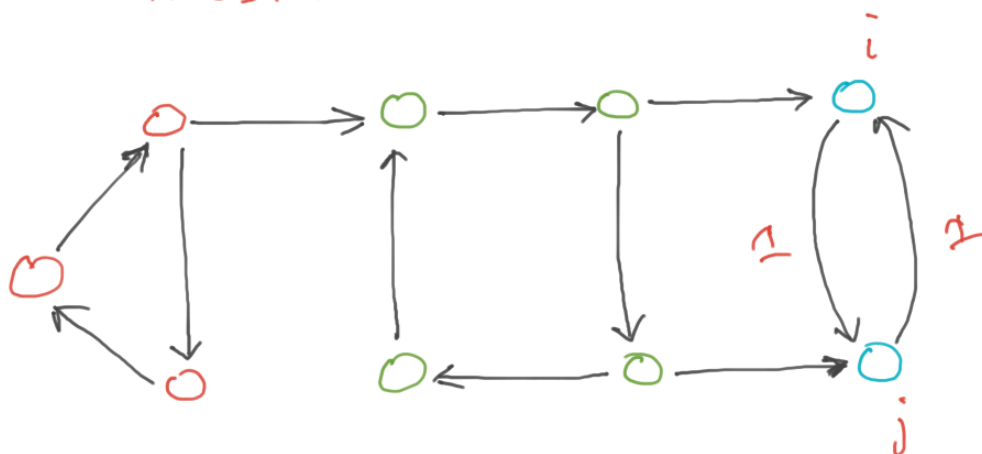
Studying distributions via "local moves"



- Often impossible to write full state space.

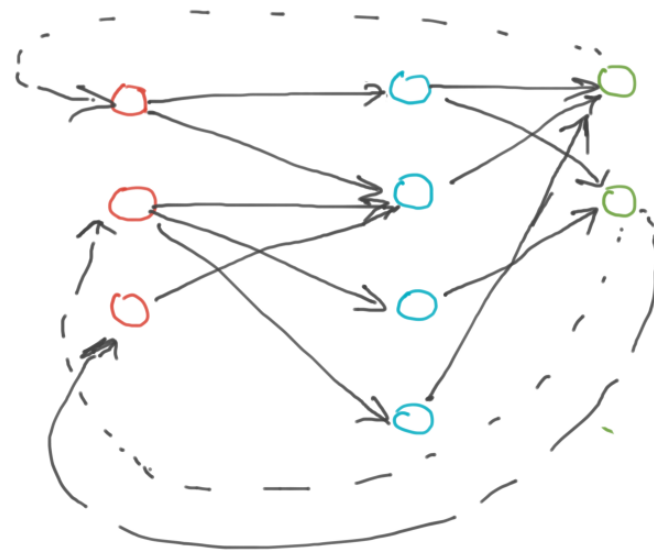
Behaviour of Markov Chains over time

Problem 1



Chain is irreducible if graph is strongly connected

Problem 2



Chain is aperiodic
if $\forall i \exists t_0$ st. $(M^t)_{ii} > 0$
 $\forall t \geq t_0$

Fundamental Thm of Markov Chains

- A finite state irreducible, aperiodic Markov chain admits a unique stationary distribution

$$\pi = (\pi_1, \dots, \pi_n) \quad (|S| = n) \quad \text{s.t.}$$

$$- \sum \pi_i = 1$$

$$- \pi M = \pi$$

$$- \forall i \quad (M^t)_{ij} \longrightarrow \pi_j \quad \text{as} \quad t \rightarrow \infty$$

A nice special case

~~Def~~ An ergodic, aperiodic Markov chain is reversible w.r.t. π if $\pi_i M_{ij} = \pi_j M_{ji}$

• If chain is reversible w.r.t. π , then

π is stationary distribution.

Speed of Convergence : Linear Algebra

For $x, y \in \mathbb{R}^n$ ($|S|=n$) define inner product

$$\langle x, y \rangle_{\pi} = \sum_{i \in S} x_i \cdot y_i \cdot \pi_i = \mathbb{E}_{i \sim \pi} [x_i \cdot y_i]$$

► M is reversible $\iff M$ is self-adjoint w.r.t.
inner-product $\langle \cdot, \cdot \rangle_{\pi}$

M is self-adjoint $\Rightarrow$ real eigenvalues $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$
 orthonormal $w_1, w_2, \dots, w_n$

converges

convergence

Ex: $|\lambda_i| \leq 1 \forall i$ (Gershgorin) and $\lambda_1 = 1$ ($M \mathbf{1} = \mathbf{1}$)

$$(M^t)_{ij} = \underbrace{(M^t e_j)}_{j^{\text{th}} \text{ col}}_i$$

$$(M^t e_j) = \sum_{k=1}^n \lambda_k^t \cdot \langle e_j, w_k \rangle \cdot w_k = \sum_{k=1}^n \lambda_k^t \cdot (w_k)_j \cdot \pi_j \cdot w_k$$

$$= \pi_j \cdot \mathbf{1} + \sum_{k=2}^n \lambda_k^t \cdot (w_k)_j \cdot \pi_j \cdot w_k$$

$$(M^t)_{ij} - \pi_j = \pi_j \cdot \sum_{k=2}^n \lambda_k^t \cdot (w_k)_j \cdot (w_k)_i$$

$$|(M^t)_{ij} - \pi_j| \leq \pi_j \cdot \sum_{k \geq 2} |\lambda_k|^t \cdot |(w_k)_i| \cdot |(w_k)_j|$$

$$\sum_j |(M^t)_{ij} - \pi_j| \leq \sum_{k \geq 2} |(w_k)_i| \cdot |\lambda_k|^t \cdot \underbrace{\sum_j \pi_j \cdot |(w_k)_j|}_{\leq 1}$$

$$\langle \mathbf{1}, (w_k)_j \rangle_{\pi} \leq \|\mathbf{1}\|_{\pi} \cdot \|w_k\|_{\pi} \leq 1$$

$$\leq \sum_{k \geq 2} \underbrace{|(w_k)_i|}_{\leq \frac{1}{\pi_i}} \cdot |\lambda_k|^t \quad (\text{why?})$$

$$\leq \frac{n}{\pi_i} \cdot |\lambda_2|^t$$

$$\therefore \sum_j |(M^t)_{ij} - \pi_j| \leq \frac{n}{\pi_i} \cdot |\lambda_2|^t$$

$$\forall i \quad \sum_j |(M^t)_{ij} - \pi_j| \leq \frac{n}{\pi_{\min}} |\lambda_2|^t = \frac{n}{\pi_{\min}} \cdot (1-\delta)^t \quad \text{if } |\lambda_2| = 1-\delta$$

Small after $t \geq \underbrace{\frac{C}{\delta} \cdot \log\left(\frac{n}{\pi_{\min}}\right)}_{\text{Mixing time}}$

- Several mathematical techniques to bound $\delta = 1 - |\lambda_2|$
- Other parameters to control mixing times

References

- Diaconis : The Markov Chain Monte Carlo Revolution
- Levin, Peres, Wilmer : Markov Chains & Mixing Times
- Blum, Hopcroft, Kannan : Foundations of Data Science (Ch. 4)

